

Design of Powered Floor Systems for Mobile Robots with Differential Evolution

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Mobile robots



Many key applications, often in bounded environments

- internal logistics
- swarm robotics

Power supply

Battery, has to be recharged ! long lasting operations?

Alternatives:

- ground + aerial (pantograph)
- wireless energy delivery
- **powered floors** and sliding contacts

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Designing a powered floor system

Choose:

- conductive/insulating strips width
- sliding contacts on robot: number and position

given some constraints

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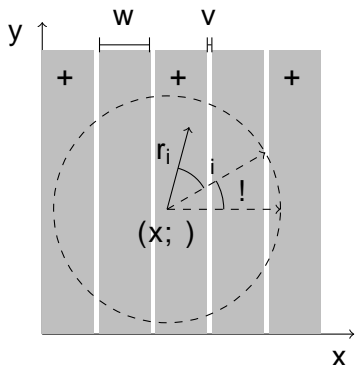
given some constraints

Can we do it automatically?

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- 3 Experiments and results

Goal



Given:

conductive (w) and insulating (v) strips width

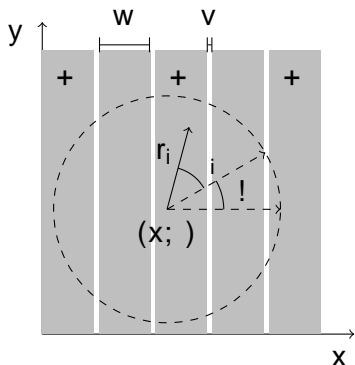
feasible region on the robot ($asg()$)

number of sliding contacts

and:

position $A = f(r_1; \alpha_1); \dots; (r_n; \alpha_n)$ of the sliding contacts array which ensure power supply in any robot position and rotation

Goal



Given:

conductive (w) and insulating (v) strips width

feasible region on the robot ($asg()$)

number of sliding contacts

and:

position $A = f(r_1; \theta_1); \dots; (r_n; \theta_n)$ of the sliding contacts array which ensure power supply in any robot position and rotation

Powered condition

"ensures power supply in any robot position and rotation"
 #
 powered condition

How to model?

one sliding contact in touch for fixed position

at least one positive and one negative in touch for fixed position

at least one positive and one negative in touch for any position

One sliding contact, $x \in \mathcal{C}$

One sliding contact; touches \mathcal{C} with robot in $x; \theta$:

$$f_c^+(r; \theta; x; \theta) = \begin{cases} 1 & \text{if } (x + r \cos(\theta)) \bmod (2w + 2v) \in [w, w + v) \\ 0 & \text{otherwise} \end{cases}$$

Similarly for $f_c^-(r; \theta; x; \theta)$

+ and , xed position

At least one positive

for xed position:

$$\exists x \in \mathbb{R}^n \left\{ \begin{array}{l} \text{touch} \\ f_c^+(r_i; i; x) \leq g(r_i; i) \end{array} \right\} \neq \emptyset \quad 1$$

+ and , fixed position

At least one positive and one negative for fixed position:

$$\begin{aligned} \exists x \in \mathbb{R}^n \{ z \mid \text{touch} \} \cap \{ z \mid \text{feasible} \} \neq \emptyset \\ f_c^+(r_i; i; x; !) g(r_i; i) \quad 1 \\ i=1 \\ \exists x \in \mathbb{R}^n \\ f_c^-(r_i; i; x; !) g(r_i; i) \quad 1 \\ i=1 \end{aligned}$$

+ and , fixed position

At least one positive and one negative for fixed position:

$$f^+(A; x; !) = \min_{z \in \mathbb{R}^n} \left\{ \begin{array}{l} \text{touch} \\ f_c^+(r_i; i; x; !) \end{array} \right\} \mid \left\{ \begin{array}{l} \text{feasible} \\ z \end{array} \right\} \quad 1$$

$$f^-(A; x; !) = \min_{z \in \mathbb{R}^n} \left\{ \begin{array}{l} f_c^-(r_i; i; x; !) \\ g(r_i; i) \end{array} \right\} \quad 1$$

Or, simply:

$$f^+(A; x; !) = \min \{ f_c^+(A; x; !); f^-(A; x; !) \} \quad 1$$

$f^+(A; x; !)$ is the number of contact pairs in touch

+ and , any position

At least one positive and one negative for any position:

$$\min_{\substack{x \in [0; 2(w+v)] \\ ! \in [0; 2]}} f^+(A; x; !)$$

+ and , any position

At least one positive and one negative for any position:

$$f(A) = \min_{\substack{x \in [0; 2(w+v)[\\ ! \in [0; 2[}} f^+(A; x; !)$$

$f(A)$ is the number of contact pairs in touch in the worst possible condition $x; !$

1 satisfies the powered condition

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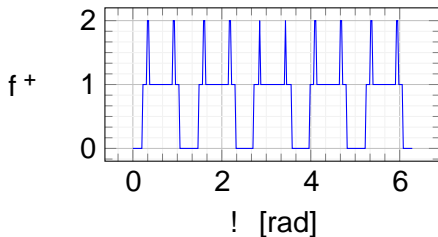
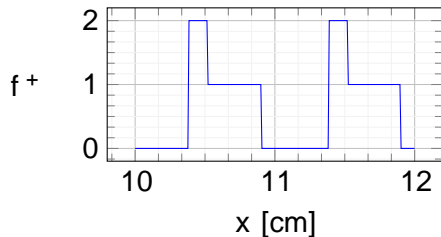
$f(A)$ is the number of contact pairs in touch in the worst possible condition $x; !$

1 satisfies the powered condition
and the larger, the better

) good candidate for being a fitness function!

What a badf !

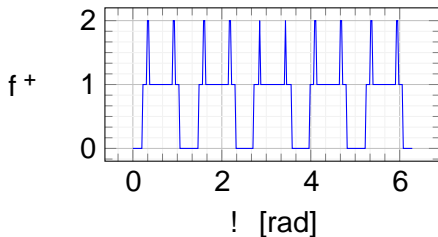
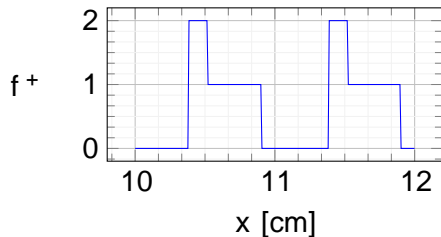
f contains a lot of bad things: min, mod, binary(!)



(for a simpleA of 5 equally spaced contacts on a circle)

What a badf !

f contains a lot of bad things: min, mod, binary(!)



(for a simpleA of 5 equally spaced contacts on a circle)

#

f might be not able to drive the search

Smootherf

Instead of \n. of contact pairs in touch in the worst condition"

$$f(A) = \min_{\substack{x \in [0; 2(w+v)] \\ \lambda \in [0; 2]}} f^+(A; x; \lambda)$$

Smootherf

Instead of "n. of contact pairs in touch in the worst condition"

$$f(A) = \min_{\substack{x \in [0; 2(w+v)] \\ ! \in [0; 2]}} f^+(A; x; !)$$

consider "averagen. of contact pairs in touch"

$$\hat{f}(A) = \frac{1}{4(w+v)} \int_{\substack{x \in [0; 2(w+v)] \\ ! \in [0; 2]}} f^+(A; x; !) \, dx \, d!$$

warning! $f = 1$ meets the powered condition, \hat{f} does not!

Practicality of an array

Are the sliding contacts “easily” doable?

Practicality of an array

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Maximize distance among contacts:

$$d(A) = \frac{1}{n} \sum_{i \in V} \min_{j \in V, j \neq i} \sqrt{r_i^2 + r_j^2 - 2r_i r_j \cos(\theta_{ij})}$$

- consider only contacts in the feasible region
- average distance to closest other contact

Recap

Given an array of sliding contacts A , we can measure:

- 1 number of contact pairs in touch in the worst condition: $f(A)$
- 2 average number of contact pairs in touch: $\hat{f}(A)$
- 3 contacts distance $d(A)$

For all: the greater, the better

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EA and solution representation

Individual:

$$\mathbf{x} = e(A) = \frac{r_1}{r}; \frac{1}{2}; \dots; \frac{r_n}{r}; \frac{n}{2}$$

- r is the largest radius in the feasible region
- $\mathbf{x} \in \mathbb{R}^{2n}$, n number of sliding contacts

Continuous optimization: we chose DE

Fitness

Three variants:

- M: only f
- MD: $f; d$
- MAD: $f; \hat{f}; d$

For MO variants, lexicographical ordering

- f comes first, it is the most important

For actual computation of f , \hat{f} , discretization of the robot position $x; !$:

- $x \in [0; 2(w + v)[$ and $! \in [0; 2 [$
- $n_{\text{point}} = 100$ each) n_{point}^2

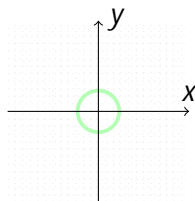
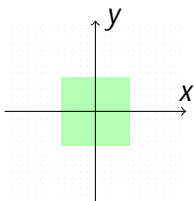
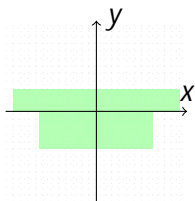
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Questions

- 1 Can f , \hat{f} , and d , possibly combined, drive the search?
- 2 Impact of parameters n , w ?

Robots: Thymio II, mBot, Elisa-3

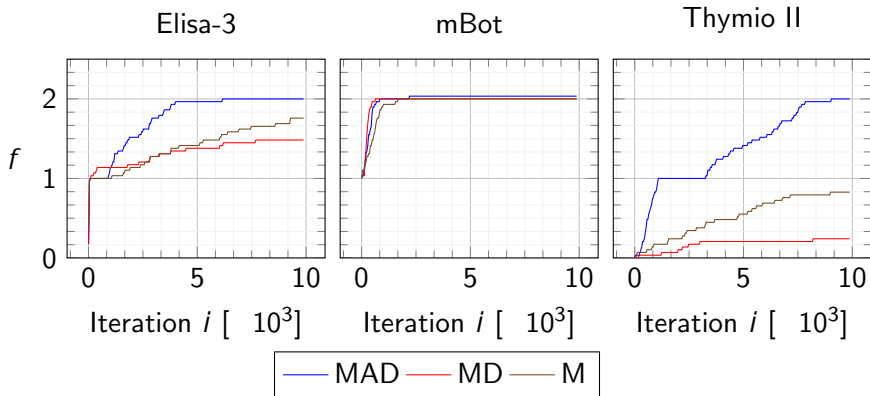


Procedure and parameters

- 30 independent runs for each problem (robot, n , w)
- 10 000 fitness evaluations per run

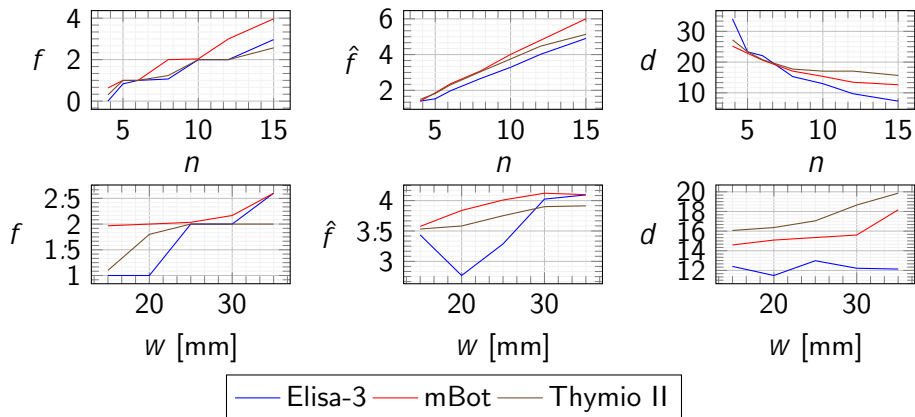
RQ1: fitness variant

$w = 25$ mm, $v = 3$ mm, and $n = 10$



- MAD is always the best
- for Thymio II, complex geometry, only MAD finds a solution!

RQ2: impact of parameters



- for n , as expected (consider cost and practicality!)
- for w , the larger, the better (v is constant)

Conclusion

Automatic design of sliding contact positions can be done!

- pretty fast: 80 s for one run
- with the proper fitness function

Future work:

- symmetry / representation
- promote balancing of + and - contacts

Thanks!